

**IN THE SPECIFICATION:**

**Kindly replace the first full paragraph on page 2 with the following:**

However, a problem arises in that if a bridge observes a net update process on one portal while processing another net update received on the other portal, two different UPDATE\_ROUTE messages related to the two different net updates can be passed to each other at the bridge without being merged into one. Then, both bridge portals are updated with the UPDATE\_ROUTE messages exchanged with each other, and then initiate bus resets on their local buses.

**Kindly replace the last paragraph on page 3 with the following:**

For example, Figure 1 illustrates a time flow diagram 100 of an example that two UPDATE\_ROUTE messages are received by a bridge on its two local buses and the two UPDATE\_ROUTE messages will be passed by and be forwarded to the other buses.

**Kindly replace the second paragraph on page 4 with the following**

The coordinator (not shown) on the first bus A (not shown) sends a net update message 120, which at point ~~120~~ 140 is received by bus ~~[[A.]]~~ B. At approximately the same time 130, the coordinator (not shown) on bus B sends a net update message to bus A, which is received at time 110. Thus there is a "crossing" of update messages between the start from one bus to the reset on the other bus, in the time frame referred to as a net update period (the period from 130 to 140).

**Kindly replace the first full paragraph on page 6 with the following:**

Since one of the net update events observed by the bridge has been discarded and only one net update event will be processed, net update events cannot be passed to each other at a bridge. That can prevent the net oscillation problem explained above.

**Kindly delete lines 10-23 on page 6.**

**Kindly delete lines 1- 24 on page 7.**

**Kindly delete lines 1 to 3 on page 8.**

**Kindly replace the last paragraph on page 8 with the following:**

With regard to how the bridge selects the survived UPDATE\_ROUTE message received by one portal and discards the other received, the following examples are for explanatory purposes only and do not limit the claimed invention to just these criteria for selecting a survived message and discarding a victim message.

- (1) First in time; the bridge keeps the first received net update message on one bus and discards the second message;
- (2) CPU priority, one CPU; the message found first by the CPU is kept and the other message is discarded; or

CPU priority multiple CPUs; the first CPU that reports a net update event detection to the other CPU ~~process~~ processes it,

**Kindly replace the first full paragraph on page 10 with the following:**

At step 200, if the portal is a coordinator on its local bus, step 210 will be processed next, otherwise if the portal is not a coordinator on its local bus, step 220 ~~420~~ will be executed next.

**Kindly replace the second full paragraph on page 10 with the following:**

At step 210, if the portal finds a net update collision on the local bus according to IEEE1394.1 bridge draft standard, step 230 430 will be next performed. Otherwise, step 200 400 will be processed next.